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Coverability and Inhibitor Arcs: an example

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1 Introduction

Coverability trees are a powerful tool in the behavioural analysis of Petri Nets. Originally introduced in [5] for Vector Addition Systems, they proved to be useful also for Place/Transition Nets (PT-nets) [3]. A coverability tree provides a finite representation of the reachable markings of a PT-net from which important behavioural properties of the net can be derived [8, 7], such as coverability of markings, boundedness, and mutual exclusion properties. Inhibitor arcs are a powerful extension of PT-nets, as they allow to check for emptiness of a place (so-called zero-testing), a feature not possible in the standard PT-net model. Actually, Place/Transition Nets with inhibitor arcs (PTI-nets) can simulate Turing machines [1] and so, many problems that are decidable for PT-nets, like reachability and boundedness, are undecidable for PTI-nets [4].

In [6], it is investigated how the standard coverability tree construction for PT-nets might be generalised to PT-nets and PTI-nets working under the a priori step semantics. This has led to a generic algorithm for the construction of so-called \textit{step coverability trees}. Even without the step semantics, inhibitor arcs introduce the problem of non-monotonicity and to deal with this, [6] first proposes a straightforward modification of the standard coverability tree construction.

As part of his BSc project, the first author investigated this modification, aimed at dealing with PT-nets with inhibitor arcs. The algorithm correctly identifies unbounded places, and always terminates in case of one inhibitor place. As shown in [6], there exists however a PT-net with three inhibitor places for which the algorithm does not terminate. It was not known whether the algorithm would always terminate in case of two inhibitor places. This paper solves this problem by giving an example of a PT-net with two inhibitor places for which the algorithm does not terminate.
2 Preliminaries

A PTI-net with inhibitor arcs (PTI-net) is a tuple \( \mathcal{N} = (P, T, W, I, M_0) \), where \( P \) and \( T \) are disjoint finite sets of places and transitions, respectively; \( W : (T \times P) \cup (P \times T) \rightarrow \mathbb{N} \) is the weight function; \( I \subseteq P \times T \) is the set of inhibitor arcs; and \( M_0 : P \rightarrow \mathbb{N} \) is the initial marking (in general, any mapping \( M : P \rightarrow \mathbb{N} \) is a marking). In diagrams places are drawn as circles, transitions as rectangles, the weight function is represented by (weighted) arcs, and inhibitor arcs are drawn with a small circle as arrowhead. A marking \( M \) is represented by placing \( M(p) \) tokens (small black dots) inside the circle representing place \( p \).

A transition \( t \in T \) is enabled at a marking \( M \) if \( M(t) \geq W(p, t) \), for every \( p \in P \), and \( M(p) = 0 \), for every \( p \in I(t) = \{ q \mid (q, t) \in I \} \) (i.e., all inhibitor places of \( t \) are empty). In such a case \( t \) can fire leading to the marking \( M' \) satisfying \( M'(p) = M(p) - W(p, t) + W(t, p) \), for every \( p \in P \). We denote this by \( M(t)M' \). The enabling condition and firing of a transition can be generalised in the obvious way to extended markings defined as mappings \( M : P \rightarrow \mathbb{N} \cup \{ \omega \} \), where \( \omega \) is the smallest infinite ordinal.

A marking \( M \) is reachable if it can be obtained from the initial marking through successively firing a finite sequence of transitions, and it is coverable if there is a reachable marking \( M' \) such that \( M'(p) \geq M(p) \), for all \( p \in P \).

A PT-net is nothing but a PTI-net without any inhibitor arcs.

3 Constructing Coverability Trees for PTI-nets

The algorithm shown in Table 1, but without the line indicated by (\( * \)), is basically the standard coverability tree construction introduced in [5]. If the input net \( N \) is a PT-net, i.e. \( I = \emptyset \), then the line indicated by (\( * \)) is void. It generates a tree \( CT = (V, A, \mu, v_0) \), where \( V \) is a set of nodes (\( v_0 \in V \) is the root), \( A \) is a set of arcs labelled with the transitions of \( N \), and \( \mu \) is a mapping associating a (possibly extended) marking with each node in \( V \). An arc labelled \( t \) from node \( v \) to node \( w \) is denoted by \( v \xrightarrow{t} w \), and \( v \sim_A w \) means that node \( w \) can be reached from node \( v \) with \( A \) being the sequence of transitions labelling the arcs along the path.

The algorithm starts with a single (root) node, corresponding to the initial marking. Then, repeatedly, for each transition that is enabled at a marking corresponding to an already generated but not yet processed node, an arc and a new node representing the resulting marking are added. If the latter already appears in the tree, the new node is ignored and otherwise it becomes an unprocessed node. For a PT-net, place unboundedness is detected whenever a marking corresponding to a new node \( M \) strictly covers the marking \( M' \) corresponding to an ancestor node (i.e., \( M' < M \) meaning that \( M' \neq M \) and \( M'(p) \leq M(p) \) for all places \( p \)). In such a case, the marking corresponding to the new node is obtained by replacing \( M(p) \) with \( \omega \) whenever \( M'(p) < M(p) \). Note that each \( \omega \) generated in this way correctly identifies place unboundedness, as the sequence of transitions between the two nodes can be repeated indefinitely starting from \( M \).

For PT-nets, the algorithm always terminates and the resulting finite coverability tree \( CT \) can be used to decide various relevant properties. In particular, one can show
Table 1. Algorithm generating a coverability tree of a PTI-net \( N = (P, T, W, I, M_0) \)

\[
CT = (V, A, \mu, v_0) \text{ where } V = \{v_0\}, A = \emptyset \text{ and } \mu(v_0) = M_0 \\
unprocessed = \{v_0\} \\
\text{while } \text{unprocessed} \neq \emptyset \\
\text{let } v \in \text{unprocessed} \\
\text{if } \mu(v) \notin \mu(V \setminus \text{unprocessed}) \text{ then} \\
\text{for every } \mu(v)[t] \mu(M) \\
V = V \cup \{w\} \text{ and } A = A \cup \{v \xrightarrow{t} w\} \\
\text{and } \text{unprocessed} = \text{unprocessed} \cup \{w\} \\
\text{if there is } u \text{ such that } u \sim_A^* v \text{ and } \mu(u) < M \\
\text{and } \mu(u)(p) < M(p) \text{ implies } I(t') = \emptyset, \text{ for all transitions } t' \text{ in } \sigma \text{ (\star)} \\
\text{then } \mu(w)(p) = (\text{if } \mu(u)(p) < M(p) \text{ then } \omega \text{ else } M(p)) \\
\text{else } \mu(w) = M \\
\text{unprocessed} = \text{unprocessed} \setminus \{v\}
\]

that all reachable markings of \( N \) are covered by the extended markings associated with the nodes of \( CT \). Moreover, for each extended marking \( M \) associated with a node of \( CT \) there are reachable markings \( M_1, M_2, \ldots \) of \( N \) approximating \( M \), i.e. for each \( n \) and every \( p \in P \), \( M_n(p) = M(p) \) if \( M(p) \in \mathbb{N} \) and \( M_n(p) \geq n \) otherwise. As an immediate consequence, \( CT \) can be used to decide unboundedness of places in reachable markings.

This standard coverability tree algorithm does not work for PTI-nets; it may in fact falsely identify places as unbounded. The problem is the non-monotonicity of PTI-nets: given two markings \( M < M' \) and a firing sequence \( \sigma \) leading from \( M \) to \( M' \), it is not guaranteed that \( \sigma \) can also be fired from \( M' \). As a consequence, the condition for generating \( \omega \)-components may be too weak. It can be strengthened by ensuring that no inhibitor features were used along the path from \( u \) to \( v \) for those places where the number of tokens has grown. This has led to the modification proposed in [6], i.e. the addition of the line marked with (\star) in the algorithm in Table 1 which is the only difference with the standard coverability tree algorithm for PT-nets: the marking corresponding to the new node is obtained by replacing each \( M(p) \) with \( \omega \) provided that \( M'(p) < M(p) \) and the transitions fired between the two nodes have no inhibitor places.

In [6] it has been shown that the algorithm in Table 1 will always terminate for a PTI-net with one inhibitor place. Moreover, an example was given that this no longer holds if the net contains three such places. The termination problem in the case of exactly two inhibitor places was left open.

4 A Counterexample

In this section we will show an example of a PTI-net with two inhibitor places for which the algorithm in Table 1 does not terminate. An important insight in the design of this
counterexample was that its two inhibitor places need to be simultaneously unbounded. This follows from the proof in [6] for PTI-nets with one inhibitor place. In short, if they would not grow simultaneously, one of the inhibitor places would generate an \( \omega \)-component. The remainder of the construction would then regard the net as a PTI-net with a single inhibitor place, and terminate.

The counterexample is shown in Figure 1. We first observe that the two inhibitor places \( p_1 \) and \( p_2 \) of \( N_0 \) are simultaneously unbounded. It is easy to see that \( N_0 \) is deterministic and can fire exactly one infinite sequence of transitions \( \sigma = \sigma_0 \sigma_1 \sigma_2 \ldots \), where \( \sigma_i = a^{2^i}cb^{2^i}d \) for every \( i \geq 0 \). For any \( \sigma_i \), the firing of transitions \( a, c \) and \( d \) does not change the total amount of tokens. However, each firing of transition \( b \) adds one token to the total count. Hence \( N_0 \) is unbounded and the inhibitor places, \( p_1 \) and \( p_2 \), are simultaneously unbounded (one only needs to consider the markings reached after firing transition sequences \( \sigma_0 \sigma_1 \ldots \sigma_i a^{2^i} \) for \( i \geq 0 \)).

**Theorem 1.** The algorithm in Table 1 does not terminate for PTI-net \( N_0 \) in Figure 1.

*Proof.* Figure 2 shows part of the reachability graph \( N_0 \) (which has the shape of an infinite line). It starts with a marking \( M_1 \) having \( x > 0 \) tokens in place \( p_1 \), and one token in place \( p_3 \). There are three situations in which markings cover ancestor markings in this fragment.

Case 1: \( M_1 = (x, 0, 1, 0) \) is covered by \( M_7 = (2x, 0, 1, 0) \) and \( M_8 = (2x - k, k, 1, 0) \) when \( 0 < k \leq x \). Between \( M_1 \) and \( M_7, M_8 \) transitions \( c \) and \( d \) fire. As \( I(c), I(d) \neq \emptyset \), the line marked with an * in the algorithm in Table 1 returns false, and so no \( \omega \) components will be produced.

Case 2: \( M_2 = (x - i, i, 1, 0) \) is covered by \( M_8 = (2x - k, k, 1, 0) \) with \( 0 < k - i \leq x \). Between \( M_2 \) and \( M_8 \) transitions \( c \) and \( d \) fire, and so no \( \omega \) components will be produced.

Case 3: \( M_3 = (0, x, 1, 0) \) is covered by \( M_8 = (2x - k, k, 1, 0) \) with \( x < k < 2x \), and by \( M_9 = (0, 2x, 1, 0) \). Between \( M_3 \) and \( M_8, M_9 \) transitions \( c \) and \( d \) fire, and so no \( \omega \) components will be produced.

We now observe that the initial marking \( M = (1, 0, 1, 0) \) is \( M_1 \) with \( x = 1 \), and that \( M_6 = M_1 \) with \( x \) replaced by \( 2x \). Consequently, the above argument can be applied to the entire infinite sequence of firings of \( N_0 \), with all possible coverings identified as in the above case analysis. As no covering leads to the generation of \( \omega \) components, the algorithm will never terminate.

\[ \square \]
\[ M_1 = (x, 0, 1, 0) \quad M_2 = (x, x - i, 1, 0) \quad M_3 = (0, x, 1, 0) \]

\[ M_4 = (0, x, 0, 1) \quad M_5 = (2j, x - j, 0, 1) \quad M_6 = (2x, 0, 0, 1) \]

\[ M_7 = (2x, 0, 1, 0) \quad M_8 = (2x - k, k, 1, 0) \quad M_9 = (0, 2x, 1, 0) \]

Fig. 2. Execution of the PTI-net in Figure 1 with \(0 < i, j < x\) and \(0 < k < 2x\).

5 Discussion

This article focused on the number of inhibitor places in a PTI-net. However, since an inhibitor place may inhibit multiple transitions, one can also focus on the number of inhibitor arcs in the PTI-net. Still, as our counterexample has only two unweighted inhibitor arcs, it also closes the gap between one inhibitor place and two inhibitor arcs. In this context it is worthwhile to notice that in [6] a construction is given which simulates the inhibitor arcs connected to a single inhibitor place by just one unweighted inhibitor arc.

We have seen that the Modified Coverability Tree Construction does not in general terminate for PTI-nets with two or more inhibitor places. This was to be expected, as the modelling power of such PTI-nets reaches Turing completeness. To have more decision power on such Petri net models, one needs to consider restricted subclasses. An example of such a subclass are the Primitive PTI-nets, introduced in [2].

References